# University of Waterloo CS240, Spring 2021 Assignment 2 Post Mortem

In general, many students left out details regarding justification of run-time, justification or correctness, etc. on Problems 1, 2, 3, and 5.

### Problem 1 [4+6=10 marks]

• For part b), some students used a linear search on the ancestors of the node at index *i*.

# Problem 2 [3+3+5+4=15 marks]

• For part b), students incorrectly interpreted the partition as running in O(n) time, resulting in a final runtime of  $O(n^2)$ .

#### Problem 3 [8 marks]

• Many students didn't provide the sorting algorithm or prove its correctness. Some others didn't find a way to sort the almost-sorted array in O(n) using swaps.

## Problem 4 [3+3+5=11 marks]

• For a), many students had O(1) for the best-case running time, which doesn't specify a lower bound (and vice versa for  $\Omega(1)$ ).

### Problem 5 [6 marks]

• Some students did not define their R when conducting radix sort, claiming that  $O(mnR) \in O(n)$  (for MSD-radix-sort).

- Students used MSD-radix-sort, which has a slower run-time than LSD-radix-sort for this problem.
- We didn't take off points for this, but some students claimed that in base n, our areas could have up to 3 digits. The upper bound for the areas is  $(n-1)^2$ , not  $n^2$

# Problem 6 [6 marks]

- Many students forgot to put a ceiling function around the final runtime. This is required for an exact lower bound.
- Some students did not use a decision tree but came up with an algorithm to find i and j by iterating through the loop. Unfortunately, this gets a lower bound of O(n) comparisons, which is not efficient enough.