




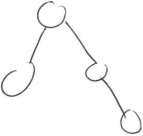


Module 4-1

4-2

AVL Trees

| | height | # nodes |
|---|--------|---------|
|  | -1 | 0 |
|  | 0 | 1 |
|  | 1 | 2 |
|   | $h=2$ | $n=3$ |
|  | $h=2$ | $n=4$ |

Height of an AVL tree

Let $N(h)$ be the least # of nodes in an AVL tree of height h

$$N(h) = \begin{cases} 1 + N(h-1) + N(h-2) & h \geq 1 \\ 1 & h=0 \\ 0 & h=-1 \end{cases}$$

Fibonacci

- $F_0 = 0$
- $F_1 = 1$
- $F_2 = 1$
- $F_3 = 2$
- $F_4 = 3$
- $F_5 = 5$

- $N(-1) = 0 = F_2 - 1$
- $N(0) = 1 = F_3 - 1$
- $N(1) = 2 = F_4 - 1$
- $N(2) = 4 = F_5 - 1$

$$N(h) = F_{h+3} - 1 = \left\lceil \frac{\varphi^{h+3}}{\sqrt{5}} \right\rceil - 1$$

golden ratio $\varphi = \frac{1 + \sqrt{5}}{2}$

solve for $h \Rightarrow \sim \log n$

$$\begin{aligned}
 N(h) &> 2N(h-2) \\
 &> 4N(h-4) \\
 &> 8N(h-6) \\
 &\vdots \\
 &> 2^i N(h-2i) \\
 &\geq 2^{\lfloor h/2 \rfloor} N(0) = 1
 \end{aligned}$$

$$n > N(h) > 2^{\lfloor h/2 \rfloor} \text{ solve for } h$$

$$\Rightarrow h \leq 2 \log n$$

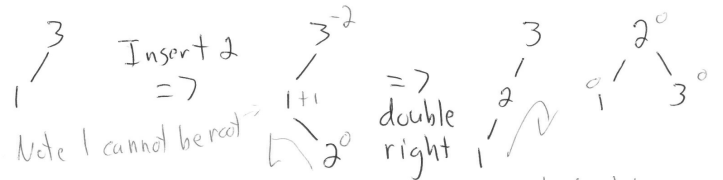
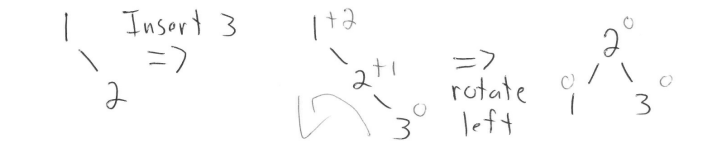
An AVL tree with n nodes has height $O(\log n)$

Both heaps & AVL tree have guaranteed height $\Theta(\log n)$. Implementations.

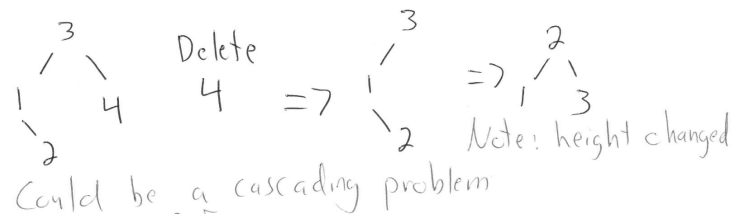
4-3

AVL - Basic Unbalanced Structures

4-4



Note: height of tree before insert is the same after -



Could be a cascading problem



Delete may require restructure up to root