## CS 240: Data Structures and Data Management

## Tutorial 8: July 4

**1.** Build a quadtree using the following points: (1, 4), (2, 5), (3, 2), (4, 7), (7, 3), (6, 1), (5, 6), (3, 7).

**2.** Build a kd-tree using the following points: (1, 4), (2, 5), (3, 2), (4, 7), (7, 3), (6, 1), (5, 6), (3, 7). Note that these are the same points as the previous problem.

**3.** Suppose that we use double hashing to resolve collisions, i.e., we use the hash function  $h(k,i) = (h_1(k) + ih_2(k)) \mod m$ . Show that if m and  $h_2(k)$  have greatest common divisor  $d \ge 1$  for some key k, then an unsuccessful insertion for key k examines  $\frac{1}{d}^{th}$  of the hash table before returning to slot  $h_1(k)$ .

Thus, when d = 1, i.e., m and  $h_2(k)$  are relatively prime, then the insertion of k can only fail if every entry of the hash table is occupied.

4. Suppose we have a hash table of size M, where M is prime and strictly greater than 2. Consider a new variation of open addressing called *Quadratic Probing*, with hash function  $h(k,i) = (h(k)+i^2) \mod M$ . Prove that Quadratic Probing will yield a maximum of only  $\lceil \frac{M+1}{2} \rceil$  distinct locations in your hash table for all  $i \in \mathbb{Z}^{\geq 0}$