

Final Practice Problem

Note: This is a sample of problems designed to help prepare for the final exam. These problems do not encompass the entire coverage of the exam (and have a heavy emphasis on post-midterm material even though the exam is cumulative), nor do they reflect the length of the exam, and should not be used as a reference for its content.

1. True or False

- (a) Final exam for CS 240 is August 10th, 9am at PAC.
- (b) A SkipList with n keys and height h must have a total of $\Theta(nh)$ nodes
- (c) In a SkipList with n keys, search operation is guaranteed to have runtime of $O(\log n)$ in the worst-case.
- (d) If we perform odd number of searches in linked list with transpose heuristic, the resulting list will always be different from initial list.
- (e) The root of compressed trie always tests the first bit
- (f) Linear probing requires two hash functions.
- (g) If α is 1 for hashing with chaining, then next `insert` is guaranteed to fail
- (h) In Cuckoo hashing, two tables can be different size
- (i) Given $h_0(k) = k \bmod 7$ with double hashing, each of size 7, $f(k) = (k \bmod 6) + 1$ is more suitable to be used as $h_1(k)$ than $g(k) = k^2 \bmod 7$
- (j) If the root of a quadtree represents the region $[0, 128) \times [0, 128)$ while deepest internal node represents the region $[88, 92) \times [24, 28)$, the height of quadtree is 6.
- (k) A range tree storing n points - where n is of the form $2^{h+1} - 1$ - such that the primary tree and all associated trees are perfect (i.e. all levels are full) has $(n+1) \log(n)$ nodes in the primary trees and all associated trees combined
- (l) A modified version of the Boyer-Moore algorithm that uses a first-occurrence array (denoting the index of the first occurrence of the argument character) instead of a last-occurrence array will always successfully find the first occurrence of a pattern P in a text T (if it appears in T).

2. Revenge of First Principle

- (a) Prove from first principles that $7n^4 - 5n^2 + 6 \in \Theta(n^4)$
- (b) Prove from first principles that $14n + 22 \in o(n \log n)$

3. Run-time Analysis

When Prashanth goes to sleep, his brain generates a number n and then runs an algorithm called $Dream(n, hour)$. This algorithm has two inputs: the number n and the hour at which Prashanth goes to sleep, which is either 9:30 PM, 10:00 PM, or 10:30 PM. The runtime of $Dream(n, hour)$ is n^i computations where i is the number of hours that Prashanth sleeps.

If Prashanth sleeps at 9:30PM, he will get between 8 to 9 hours of sleep. If he goes to sleep at 10:00 PM, he will get exactly 7.5 hours of sleep. If he sleeps at 10:30 PM, he will get between 5 to 7 hours of sleep.

Prove or disprove the following statements:

- (a) In the worst case, the running time of $Dream$ is $\Theta(n^9)$.

- (b) In the best case, the running time of *Dream* is $\omega(n^5)$
- (c) If Prashanth sleeps at 10:00 PM, the running time of *Dream* is $\Omega(n^{7.5})$
- (d) The running time of *Dream* is $O(i^n)$.
- (e) After learning about some data structures in CS 240 that could help speed up his algorithm, Prashanth speeds up his *Dream* algorithm so that it runs in $\Theta(i)$ time where i is still the number of hours that he sleeps. The running time of *Dream* is $o(1)$.

4. Partial Sum Part 2

Consider the problem where we have a sequence of n elements: $S = a_1, a_2, \dots, a_n$, and 3 operations:

- $Add(S, b) \rightarrow a_1, a_2, \dots, a_n, b$
- $Update(S, i, \Delta) \rightarrow a_1, \dots, a_{i-1}, \Delta, a_{i+1}, \dots, a_n$
- $PartialSum(S, k) \rightarrow \sum_{i=1}^k a_i$

Design a data structure that can perform each of these operations in $O(\log n)$ expected time.

NOTE: we did solve this problem in tutorial 07 using AVL-tree. So, do NOT use AVL-tree this time.

5. Find Largest x in kd-tree

Given a 2D kd-tree with n points, design an algorithm that finds a point with largest x -coordinate. For simplicity, you may assume n is a power of 4.

6. Consecutive Trie Strings

Given an uncompressed trie T that stores a list of binary strings, design an algorithm $Consecutive(b_1, b_2)$ that takes two binary strings in T as input, and outputs true if the strings are consecutive in pre-order traversal of the trie, and outputs false otherwise. Assume that branches are ordered as 0, 1. The runtime should be bounded by $O(|b_1| + |b_2|)$.

For example, suppose T stores $\{000, 01, 0110, 101, 11\}$.

$Consecutive(0110, 101)$ outputs true.

$Consecutive(01, 000)$ outputs true.

$Consecutive(11, 000)$ outputs false.

7. MaxDiff

Consider an array A of n integers. We want to implement a range query called $MaxDiff(i, j)$ which will find the maximal difference between two elements from $A[i]$ to $A[j]$ inclusive, for $i < j$. For example, suppose our array A is:

$A = 3\ 0\ 5\ 4\ 5\ 6\ 3\ 4\ 5\ 7\ 9\ 8\ 1\ 0\ 1$

If we run the query $MaxDiff(2, 9)$, then the subarray from indices 2 to 9 is:

$A[2 \dots 9] = 5\ 4\ 5\ 6\ 3\ 4\ 5\ 7$

The largest number is 7 and the smallest number is 3, so the maximal difference is $7 - 3 = 4$. The query $MaxDiff(2, 9)$ should return 4.

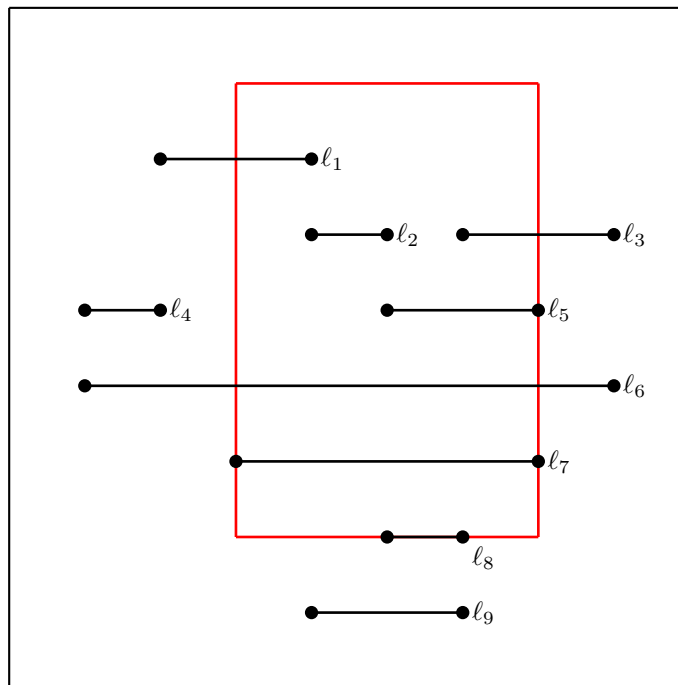
Design a data structure for A with space complexity $O(n)$ to answer queries of the form $MaxDiff(i, j)$ in $O(\log n)$ time. There are no constraints on the runtime for preprocessing the array into the data structure.

8. Range Tree

Suppose you have a set of n horizontal line segments in a plane, where line segment ℓ_i has coordinates (x_i, y_i) and (x'_i, y_i) . Assume that all coordinates are integers.

For each of the range-search queries below, design a data structure and provide an algorithm to answer the queries in $O(\log^3 n + s)$ time, where s is the number of lines reported. Each range-search query is a rectangle of the form $[a, b] \times [c, d]$.

- The algorithm reports all line segments that are entirely contained inside the query rectangle. For the example below, the algorithm would return ℓ_2, ℓ_5, ℓ_7 and ℓ_8 .
- The algorithm reports all line segments that intersect the query rectangle. For the example below, the algorithm reports all line segments **except** ℓ_4 and ℓ_9 .



9. Hashing

Let $p \geq 3$ be prime, and consider the universe of keys $U = \{0, 1, \dots, p^2 - 1\}$.

- With a hash table of size p , and using double hashing with $h_0(k) = k \bmod p$ and $h_1(k) = \lfloor k/p \rfloor + 1$, give a sequence of **two** keys to be inserted that results in failure.
- With two hash tables of sizes p and $(p - 1)$, and using cuckoo hashing with $h_0(k) = k \bmod p$ and $h_1(k) = k \bmod (p - 1)$, give a sequence of **four** keys to be inserted that results in failure.

10. Chaotic Karp-Rabin

Consider following hash function to be used in Karp-Rabin algorithm.

$$h(P) = \text{Number of vowels in } P$$

- Describe how fast-update on fingerprint must be done.
- Find an example pattern and text that leads to a case that maximizes execution of `strcmp`.

11. **Basics of Range Tree** Consider the following points being stored in a 2D range tree: $(2, 12)$, $(17, 77)$, $(23, 92)$, $(40, 47)$, $(55, 91)$, $(67, 27)$, $(89, 79)$, $(99, 53)$, $(10, 23)$, $(35, 7)$, $(61, 40)$, $(95, 56)$, $(22, 42)$, $(88, 15)$, $(42, 2)$.
- (a) Draw the primary tree for this range tree.
 - (b) Draw the corresponding associate trees for the points $(88, 15)$, $(61, 40)$ and $(67, 27)$.
 - (c) Perform a range-search with the query rectangle $[35, 90] \times [5, 30]$, indicating all boundary nodes and topmost inside nodes.
12. **Modified AVL Tree**
We consider a modified version of AVL trees where the height difference between the right and left subtrees of any node is in the range $[-2, 2]$ instead of $[-1, 1]$. These are called AVL-2 trees. Prove that the height of an AVL-2 tree on n nodes is in $O(\log n)$.
13. **Tracing SkipList**
Insert the numbers 12, 11, 13, 10, 20 into an empty skip-list using the coin flips HHTHTHTTHHHT. Then delete the keys 13 and 20.
14. **Boyer More**
Let the text be "BATARATATATCAT", and the pattern be "TATCAT", show the full Boyer Moore Algorithm applied to this string search.