• Data definition and templates
• Subintervals
• Counting up

Natural Numbers
Readings

• HtDP, sections 11 – 13 (Intermezzo 2)
Data definition and templates
Natural numbers

;; A Nat is one of:
;; * 0
;; * (add1 Nat)

Here, add1 is the built-in function that adds 1 to its argument.

The natural numbers start at 0 in computer science and some branches of mathematics (e.g., logic).

We will now work out a template for functions that consume a natural number.
Suppose we have a natural number n.

;; A Nat is one of:
;; 0
;; (add1 Nat)

The test (zero? n) tells us which case applies.
If (zero? n) is false, then n has the value (add1 k) for some k.
Suppose we have a natural number n.

;; A Nat is one of:
;; * 0
;; * (add1 Nat)

To compute k, we subtract 1 from n, using the built-in sub1 function.
Because the result (add1 n) is a natural number, we recursively apply the function we are constructing to it.

(define (nat-template n)
  (cond
   [(zero? n) ...]
   [else (...[...n
   (nat-template(sub1 n)))])))
Example: a decreasing list

Countdown function: it consumes a natural number $n$ and produces a decreasing list of all natural numbers less than or equal to $n$.

(countdown 0) => (cons 0 empty)
(countdown 2) => (cons 2 (cons 1 (cons 0 empty)))

With these examples, we proceed by filling in the template.
Example: a decreasing list

;; (countdown n) produces a list of decreasing Nats
;; from n to 0

;; countdown: Nat -> (listof Nat)

(check-expect (countdown 0) (cons 0 empty))

(check-expect (countdown 2)
  
  (cons 2 (cons 1 (cons 0 empty)))
)

(define (countdown n)
  (cond
    [(zero? n) ...]
    [else (...]
      ...
      n
      (countdown (sub1 n)))))

If $n = 0$, we produce the list containing 0, and if $n \neq 0$, we cons $n$ onto the countdown list for $n - 1$. 
Example: a decreasing list

;; (countdown n) produces a list of decreasing Nats
;; from n to 0
;; countdown: Nat -> (listof Nat)
(check-expect (countdown 0) (cons 0 empty))
(check-expect (countdown 2)
             (cons 2 (cons 1 (cons 0 empty))))
(define (countdown n)
  (cond
   [(zero? n) (cons 0 empty)]
   [else (cons
           n
           (countdown (sub1 n)))]))
A condensed trace

(define (countdown n)
  (cond
    [(zero? n) (cons 0 empty)]
    [else (cons n
      (countdown (sub1 n)))]
  ))

(countdown 2)
=> (cons 2 (countdown 1))
=> (cons 2 (cons 1 (countdown 0)))
=> (cons 2 (cons 1 (cons 0 empty)))
Subintervals
Subintervals of the natural numbers

The symbol $\mathbb{Z}$ is often used to denote the integers. We can add subscripts to define subsets of the integers.

For example, $\mathbb{Z}_{\geq 0}$ defines the non-negative integers, also known as the natural numbers.

Other examples include $\mathbb{Z}_{>4}$, $\mathbb{Z}_{<-8}$, and $\mathbb{Z}_{\leq 1}$.
Example: $\mathbb{Z}_{\geq 7}$

If we change the base case test from `(zero? n)` to `(= n 7)`, we can stop the countdown at 7.

This corresponds to the following definition:

```plaintext
;; An integer in Z>=7 is one of:
;; * 7
;; * (add 1 Z>=7)
```

We use this data definition as a guide when writing functions, but in practice we use a requires section in the contact to capture the new stopping point.
;; (countdown-to-7 n) produces a decreasing list from n to 7
;; countdown-to-7: Nat -> (listof Nat)
;; requires: n >= 7
;; Example:
(check-expect (countdown-to-7 9)
  (cons 9 (cons 8 (cons 7 empty)))))

(define (countdown-to-7 n)
  (cond
    [ (= n 7) (cons 7 empty)]
    [else (cons n (countdown-to-7 (sub1 n)))]))
Generalizing **countdown** and **countdown-to-7**

We can generalize both **countdown** and **countdown-to-7** by providing the base value (e.g., 0 or 7) as a second parameter \( b \) (the "base").

Here, the stopping condition will depend on \( b \).

The parameter \( b \) has to "go along for the ride" (be passed unchanged) in the recursion.
;; (countdown-to n b) produces a decreasing list from
;; n to b
;; countdown-to: Int Int -> (listof Int)
;; requires: n >= b
;; Example:
(check-expect (countdown-to 4 2)
  (cons 4 (cons 3 (cons 2 empty))))
(define (countdown-to n b)
  (cond
   [ (= n b) (cons b empty)]
   [else (cons n (countdown-to (sub1 n) b)]))
Another condensed trace

\[
\text{(define} \ (\text{countdown-to} \ n \ b) \\
\text{(cond} \\
\quad [(= n b) (\text{cons} \ b \ \text{empty})] \\
\quad [\text{else} \ (\text{cons} \ n \ (\text{countdown-to} \ (\text{sub1} \ n) \ b))])]))) \\
\text{(countdown-to} \ 4 \ 2) \\
\Rightarrow (\text{cons} \ 4 \ (\text{countdown-to} \ 3 \ 2)) \\
\Rightarrow (\text{cons} \ 4 \ (\text{cons} \ 3 \ (\text{countdown-to} \ 2 \ 2))) \\
\Rightarrow (\text{cons} \ 4 \ (\text{cons} \ 3 \ (\text{cons} \ 2 \ \text{empty})))
\]
countdown-to with negative numbers

(define (countdown-to n b)
  (cond
    [(= n b) (cons b empty)]
    [else (cons n (countdown-to (sub1 n) b))]))

(countdown-to 1 -2)
=> (cons 1 (countdown-to 0 -2))
=> (cons 1 (cons 0 (countdown-to -1 -2))
=> (cons 1 (cons 0 (cons -1 (countdown-to -2 -2))))
=> (cons 1 (cons 0 (cons -1 (cons -2 empty)))))
Counting up
Counting up

What if we want an increasing count?
Consider the non-positive integers $\mathbb{Z}_{\leq 0}$.

;;; A integer in $\mathbb{Z} \leq 0$ is one of:
;;; * 0
;;; * (sub1 $\mathbb{Z} \leq 0$)

Examples:
• $-1$ is (sub1 0)
• $-2$ is (sub1 (sub1 0))

If an integer $i$ is of the form (sub1 $k$), then $k$ is equal to (add1 $i$).
This suggests the following template.
Notice the additional requires section.

;; nonpos-template: Int -> Any
;; requires: n <= 0
(define (nonpos-template n)
  (cond
   [(zero? n) ...]
   [else (...
     ...n
     (nonpos-template (add1 n)))]))

We can use this to develop a function to produce lists such as
(cons -2 (cons -1 (cons 0 empty))).
;; (countup n) produces an increasing list from n to 0
;; countup: Int -> (listof Int)
;; requires: n <= 0

;; Example:
(check-expect (countup -2) (cons -2 (cons -1 (cons 0 empty)))))

(define (countup n)
  (cond
   [(zero? n) (cons 0 empty)]
   [else (cons n (countup (add1 n)))]))
As before, we can generalize this to counting up to \( b \), by introducing \( b \) as a second parameter in a template.

\[
\text{(define \ (upto-b-template \ n \ b) \}
\begin{aligned}
&\text{(cond} \\
&\quad \text{[ (= n b) \ (b ...)]} \\
&\quad \text{[else \ ( ...} \text{n} \\
&\quad\quad \text{(upto-b-template \ (add1 n) b))])})
\end{aligned}
\]
;; (countup-to n b) produces an increasing list from n
to b
;; countup-to: Int Int -> (listof Int)
;; requires: n ≤ b
;; Example:
(check-expect (countup-to 6 8)
  (cons 6 (cons 7 (cons 8 empty))))
(define (countup-to n b)
  (cond
    [(= n b) (cons b b empty)]
    [else (cons n (countup-to (add1 n) b))])))
Many imperative programming languages offer several language constructs to do repetition, e.g.,

```plaintext
for i = 1 to 10 do { ... }
```

Racket offers one construct – recursion – that is flexible enough to handle these situations and more.

We will soon see how to use Racket’s abstraction capabilities to abbreviate many common uses of recursion.
When you are learning to use recursion, sometimes you will “get it backwards” and use the countdown pattern when you should be using the countup pattern, or vice-versa.

Avoid using the built-in list function `reverse` to fix your error. It cannot always save a computation done in the wrong order.

Instead, learn to fix your mistake by using the right pattern.

You may not use `reverse` on assignments unless we say otherwise!
End of module
Goals of this module

• You should understand the recursive definition of a natural number, and how it leads to a template for recursive functions that consume natural numbers.

• You should understand how subsets of the integers greater than or equal to some bound \( m \), or less than or equal to such a bound, can be defined recursively, and how this leads to a template for recursive functions that “count down” or “count up”. You should be able to write such functions.