Error detection

1. The variable ‘y’ is declared twice. This is a semantic error.
2. The variable ‘idx’ is not declared at all. This is a semantic error.
3. In WLP4, all variable declarations must proceed all statements. This is a syntax error, since the WLP4 grammar forces this structure.
4. In WLP4, all procedure declarations must also be procedure definition. This is a syntax error.
5. Cannot assign the output of a function that returns a float to a variable of type int. This is a semantic error.
6. The function ‘getRandom()’ is not declared. This is a semantic error.

Expression code generation

Immediately, we must first recall what these two operators do. Pre-increment will increment the value of ID by 1 and return that new value. Post-increment will also increment the value of ID by 1 but will return the old value.

Pre-increment is easier so let’s do that first.

```c
void genCode(tree t)
{
    ...
    if(t.rule is "factor -> PLUS PLUS lvalue"){
        // generate code to put the lvalue in $3
        // You’ll do this in your assignment
        genCode(t.children[2])
        // $3 contains an lvalue, which means we need to fetch the actual value
        // from memory
        // We're making the assumption that the full address and not just an offset
        // Otherwise we need to do something like: add/sub $3, $3, $29
        lw $5, 0($3)
        // Add one to the value
        add $5, $5, $11
        // Save the new value
        sw $5, 0($3)
        // Return new value
        add $3, $5, $0
    }
}
Note that we return the actual value and not an lvalue since we have no rules that reduce factor to lvalue (e.g. \texttt{factor -> lvalue}). This also means we cannot nest pre and post-increment (e.g. \texttt{++(++i)}). Similarly, we cannot return an lvalue as we have no grammar rules to support \texttt{expr -> expr + lvalue}. We could add additional rules to facilitate this but that is left as an exercise to you.

Post-increment is “harder” only in the sense that we must keep a copy of the old value and return that instead.

```c
void genCode(tree t)
{
    if(t.rule is "factor -> lvalue PLUS PLUS"){
        // generate code to put the lvalue in $3
        // You’ll do this in your assignment
        genCode(t.children[0])
        // $3 contains an lvalue, which means we need to fetch the actual value
        // from memory
        // We’re making the assumption that the full address and not just an offset
        // Otherwise we need to do something like: add/sub $3, $3, $29
        lw $5, 0($3)
        // Copy the old value
        add $6, $5, $0
        // Add one to the current value
        add $5, $5, $11
        // Save the new value
        sw $5, 0($3)
        // Return old value
        add $3, $6, $0
    }
}
```

As we can see, post-increment is not that much more expensive than pre-increment for integers. It’s one instruction more but this gets compounded as we create more complex data structures.

Switch statement

Here the individual case statements need to be able to jump to a label that the switch statement generated, so we need a way to pass the label IDs down the tree. We will do this by creating a \texttt{parentLabelId} field which the parent passed down the tree to its children.

```c
void genCode(tree t)
{
    if(t.rule is "statement -> SWITCH LPAREN expr RPAREN LBRACE cases default RBRACE"){
        x = genLabelID()
        // Evaluate the expr and push it onto the stack
        genCode(t.children[2])
        push($3)
        // Pass the label ID we generated to the children
        c.children[5].parentLabelId = x
        // Generate all the case statements
    }
}
```
Test code generation

Our immediate instinct to generate code for the less than or equal to test is:

- Generate code for left expression and push result onto stack
- Generate code for right expression and pop previous result from stack into $5.
  - $3 contains right-value, $5 contains left-value
- Check if $5 is less than $3
- If it is, return 1
- Otherwise, check if they are equal
- If they are equal, return 1
- Otherwise, return 0

What is the problem with this code? Well, it’s slightly inefficient. We use more instructions than necessary. Let’s see:
slt $6, $5, $3
bne $0, $6, returnOne
beq $5, $3, returnOne
add $3, $0, $0
beq $0, $0, end
returnOne:
add $3, $11, $0

Is there a better way? Yes. What does $X \leq Y$ mean? It evaluates to true if X is not greater than Y. Equivalently, it returns true if Y is not less than X, i.e. !(Y < X). But that’s effectively swapping the arguments in our slt expression.

slt $6, $3, $5
sub $3, $11, $6 ; if $6 = 1$ then $3 = 0 \Rightarrow y < x$

So we’ve managed to significantly reduce our generated code from our initial version by the use of some fancy negation techniques.