Lecture 2: MergeSort

CS 341: Algorithms

Thu, Jan 10th 2019
Outline For Today

1. Example 1: Sorting-Merge Sort-Divide & Conquer
Sorting

◆ Input: An array of integers in *arbitrary* order

10  2  37  5  9  55  20

◆ Output: Same array of integers in *increasing* order

2  5  9  10  20  37  55
Algorithm 1: Selection Sort

procedure selectionSort(Array X of size n):
    for i = 1 to n {
        let minIndex = i;
        for j = i+1 to n {
            if X[j] < X[minIndex]
                minIndex = j
        }
        X[i] <-- X[minIndex] (swap in place)
    }
    return X
SelectionSort Simulation

i = 1

minElement: 2
minIndex: 2
SelectionSort Simulation

minElement: 5
minIndex: 4
SelectionSort Simulation

minElement: 9
minIndex: 5

i = 3
SelectionSort Simulation

minElement: 10
minIndex: 4
SelectionSort Simulation

2  5  9  10  37  55  20

i=5

minElement: 20
minIndex: 7
SelectionSort Simulation

minElement: 37
minIndex: 7
SelectionSort Simulation

Final Output

| 2 | 5 | 9 | 10 | 20 | 37 | 55 |
procedure selectionSort(Array X of size n):
    for i = 1 to n {
        let minIndex = i;
        for j = i+1 to n {
            if X[j] < X[minIndex]
                minIndex = j
        }
        X[i] <-> X[minIndex]
    }
return X
Analysis of Selection Sort

◆ Inner Loop: $3(n-1) + 3(n-2) + \ldots + 3$

$$3 \sum_{k=1}^{n-1} k = \frac{3(n-1)n}{2} = \frac{3n^2 - 3n}{2}$$

◆ Outer Loop: $n$

◆ Initial Assignment: $n$

◆ Swap: $n$

◆ Total: $\frac{3n^2 + 3n}{2}$

**SelectionSort takes $(3n^2 + 3n)/2$ time on an input of size $n$.**
Criticism of Our Analysis & Sloppiness in CS 341

◆ Criticism 1: Loop increment is not 1 but 2 operations.
◆ Criticism 2: Swap is not 1 but 3 operations.
◆ Criticism 3: At machine level, swap might be 100 operations.

In CS 341, we’ll be sloppy in our counting of what constitutes how many operations.

We’ll count as “1 operation” high-level operations such as addition/subtraction/comparison/swap, etc.

◆ Will make more formal with Big-oh notation in a few lectures
Algorithm 2: MergeSort (Divide & Conquer)

- Assume n is power of 2. (Doesn’t really matter)

**procedure** mergeSort(Array X of size n):

1. mergeSort(left subarray X[1,...,n/2])
2. mergeSort(right subarray X[n/2+1,...,n])
3. combine the left & right sorted halves

Will simulate how the third step is done.
MergeSort Downward-Recursive Steps
MergeSort Upward-Recursive Steps
Merge Subroutine Simulation 1

L

5 31

R

12 21

O

5 12 21 31
Combine Subroutine Simulation 2

L

4 10 96 98

R

5 12 21 31

O

4 5 10 12 21 31 96 98
Pseudocode for Merge Subroutine

procedure merge(sorted lists L,R of size m/2):
    Out = empty array of size m
    i = 1; j = 1;
    for k = 1 to m:
        if L[i] < R[j]:
            Out[k] = L[i];
            i++;
        else:
            Out[k] = R[j];
            j++;
Analysis of MergeSort

Q: How much time (# operations) does MergeSort take?

**procedure** mergeSort(Array X of size n):

1. mergeSort(left subarray X[1, ..., n/2])
2. mergeSort(right subarray X[n/2+1, ..., n])
3. merge the left & right sorted halves

Simpler Question: How many operations does the merge subroutine (step 3 take) on an input of size m?
Analysis for Merge Subroutine

**procedure** merge(sorted lists L,R of size m/2):
  Out = empty array of size m → m ops
  i = 1; j = 1; → 2 ops
  **for** k = 1 to m: → 1 op
    **if** L[i] < R[j]: → 1 op
      Out[k] = L[i]; → 1 op
      i++; → 1 Op
    else:
      Out[k] = R[j]; → 1 op
      j++; → 1 Op
  → 4 ops

Total: m + 2 + 4m = 5m + 2 ≤ 7m
**Analysis of MergeSort**

**MergeSort takes** \(7n \log_2(n) + 7n\) time on an input of size \(n\).
Fundamental (\& Fast) Algorithms to Tractable Problems

- MergeSort
  - Strassen's MM
  - BFS/DFS
  - Dijkstra's SSSP
  - Kosaraju's SCC
  - Kruskal's MST
  - Floyd Warshall APSP
  - Topological Sort
  - ...

Common Algorithm Design Paradigms

- Divide-and-Conquer
- Greedy
- Dynamic Programming

Mathematical Tools to Analyze Algorithms

- Big-Oh notation
- Recursion Tree
- Master method
- Substitution method
- Exchange Arguments

Intractable Problems

- P vs NP
- Poly-time Reductions
- Undecidability

Other (Last Lecture)

- Randomized/Online/Parallel Algorithms
1. Worst-case Runtime Analysis
   - Justification 1: Easier to make worst-case analysis
   - Justification 2: Holds under any input => Very strong statement

2. “Sloppy” in counting
   - Justification: Can agree on high-level ops but impossible to agree on low-level ops
     - Will be different from language to language/architecture to architecture/compiler to compiler
3. Interested in very large inputs

- Mathematically can’t say $3n^2 + 3n/2 > 7n\log_2(n) \Rightarrow$ depends on $n$
- But for large $n$, can say that $3n^2 + 3n/2 > 7n\log_2(n)$
- So we’ll say: MergeSort is better than SelectionSort