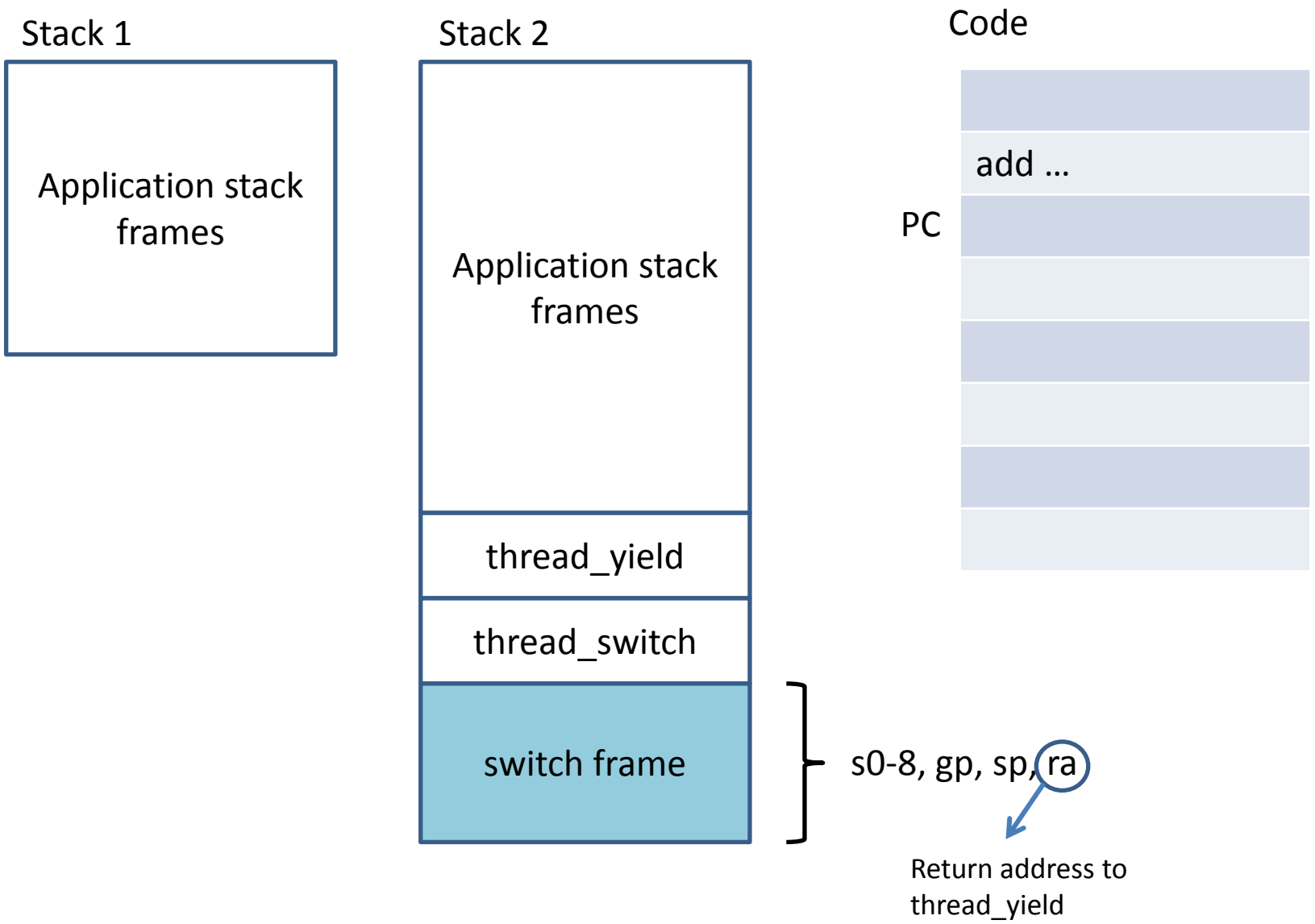
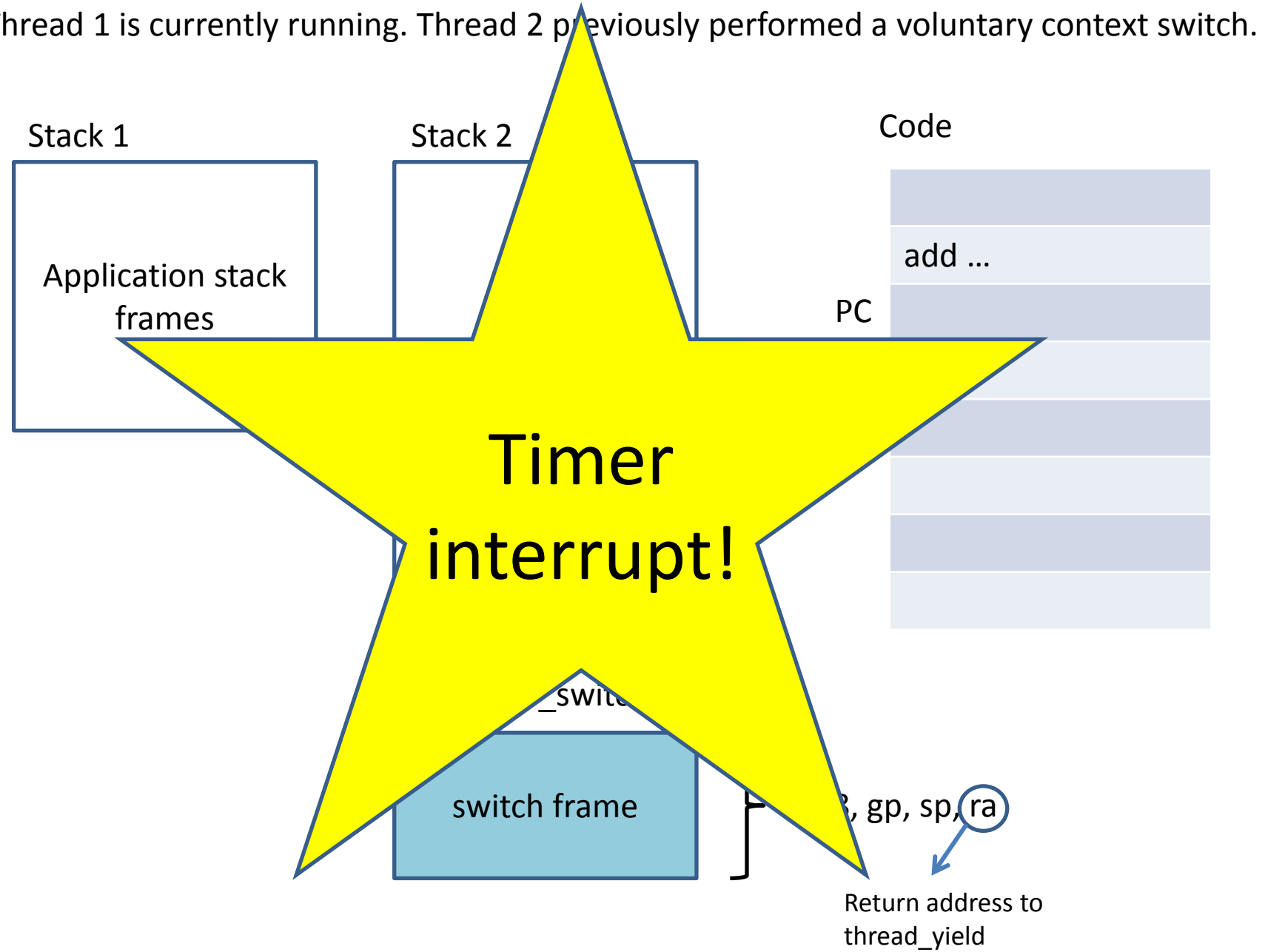
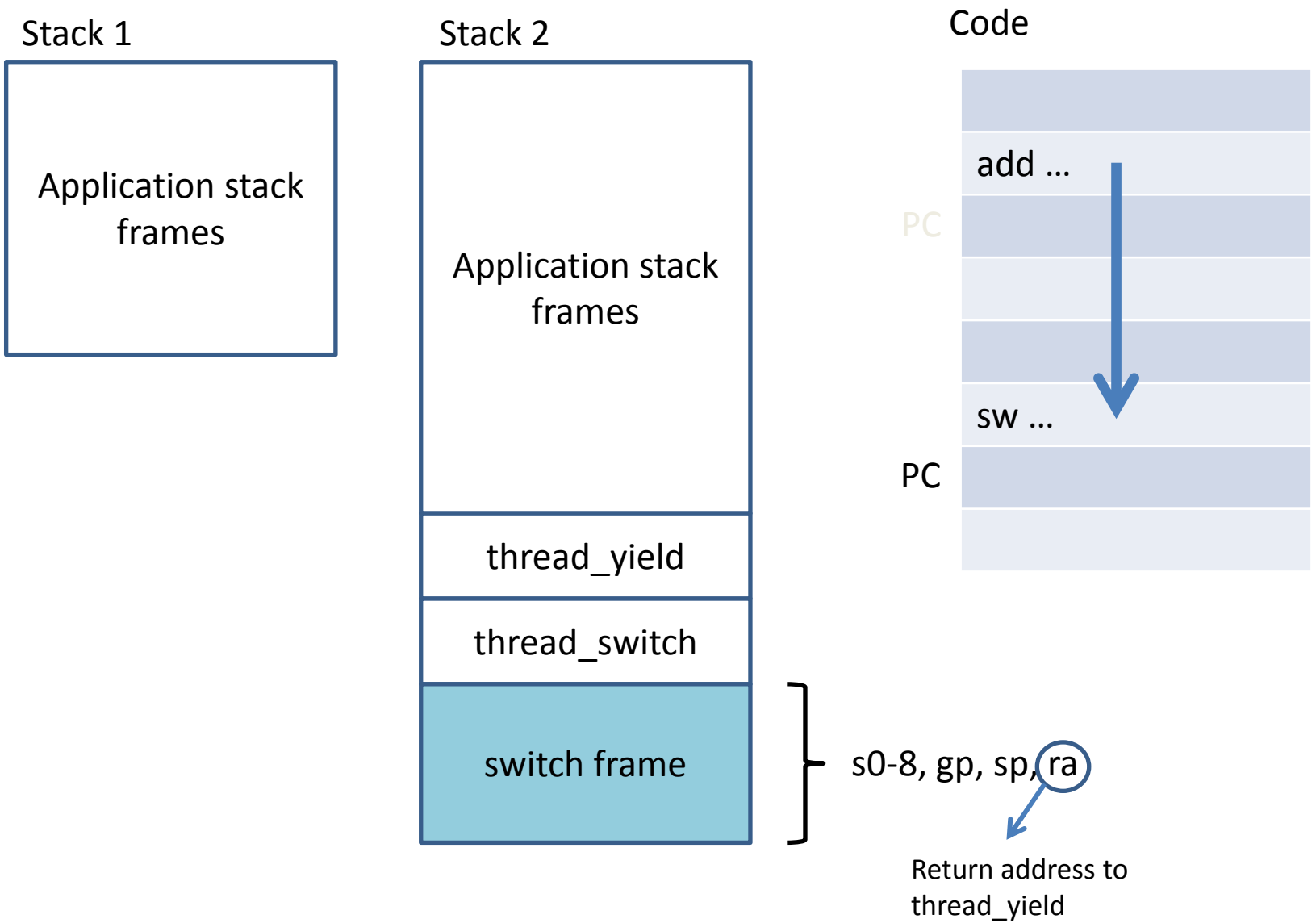
Thread 1 is currently running. Thread 2 previously performed a voluntary context switch.



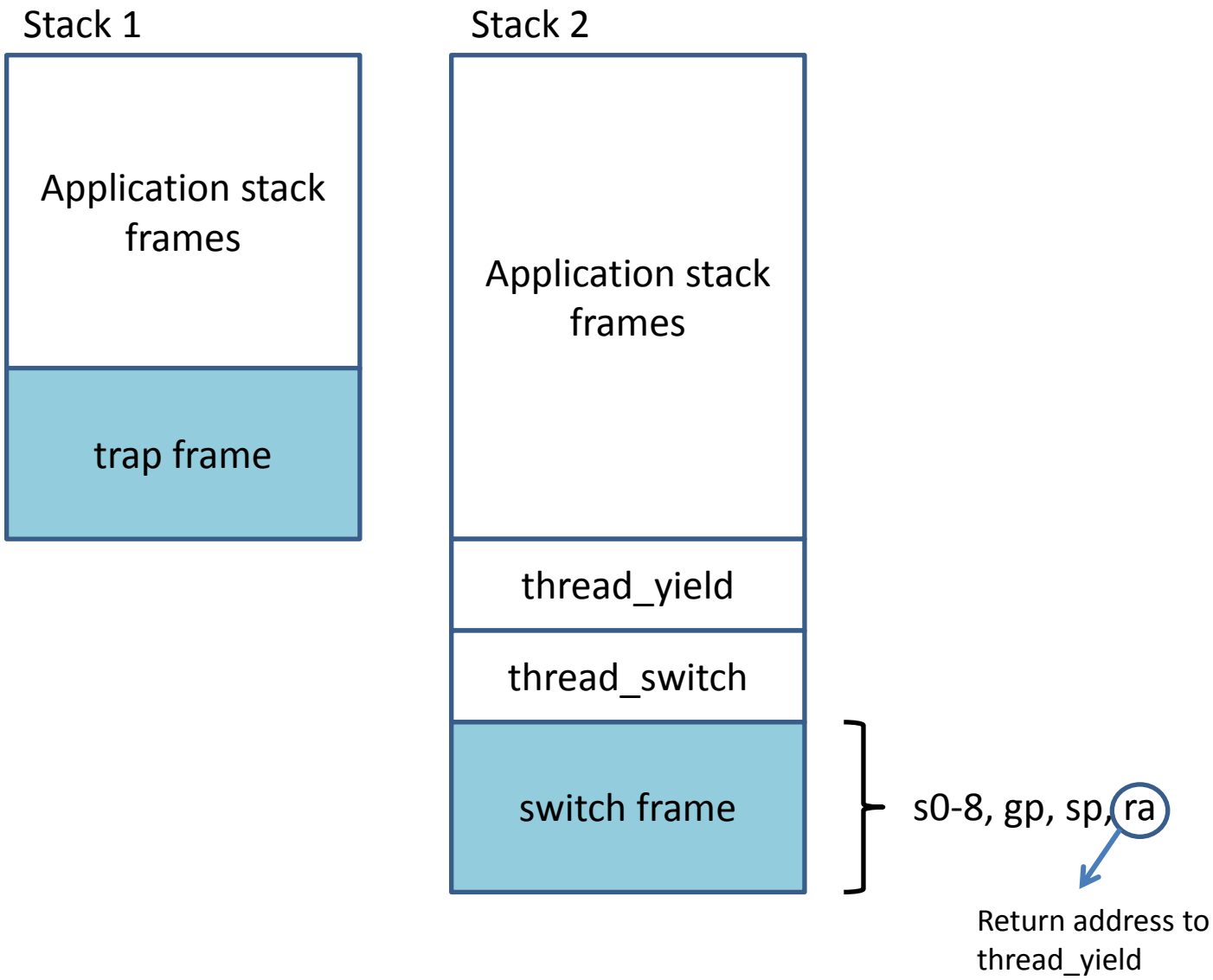
Thread 1 is currently running. Thread 2 previously performed a voluntary context switch.



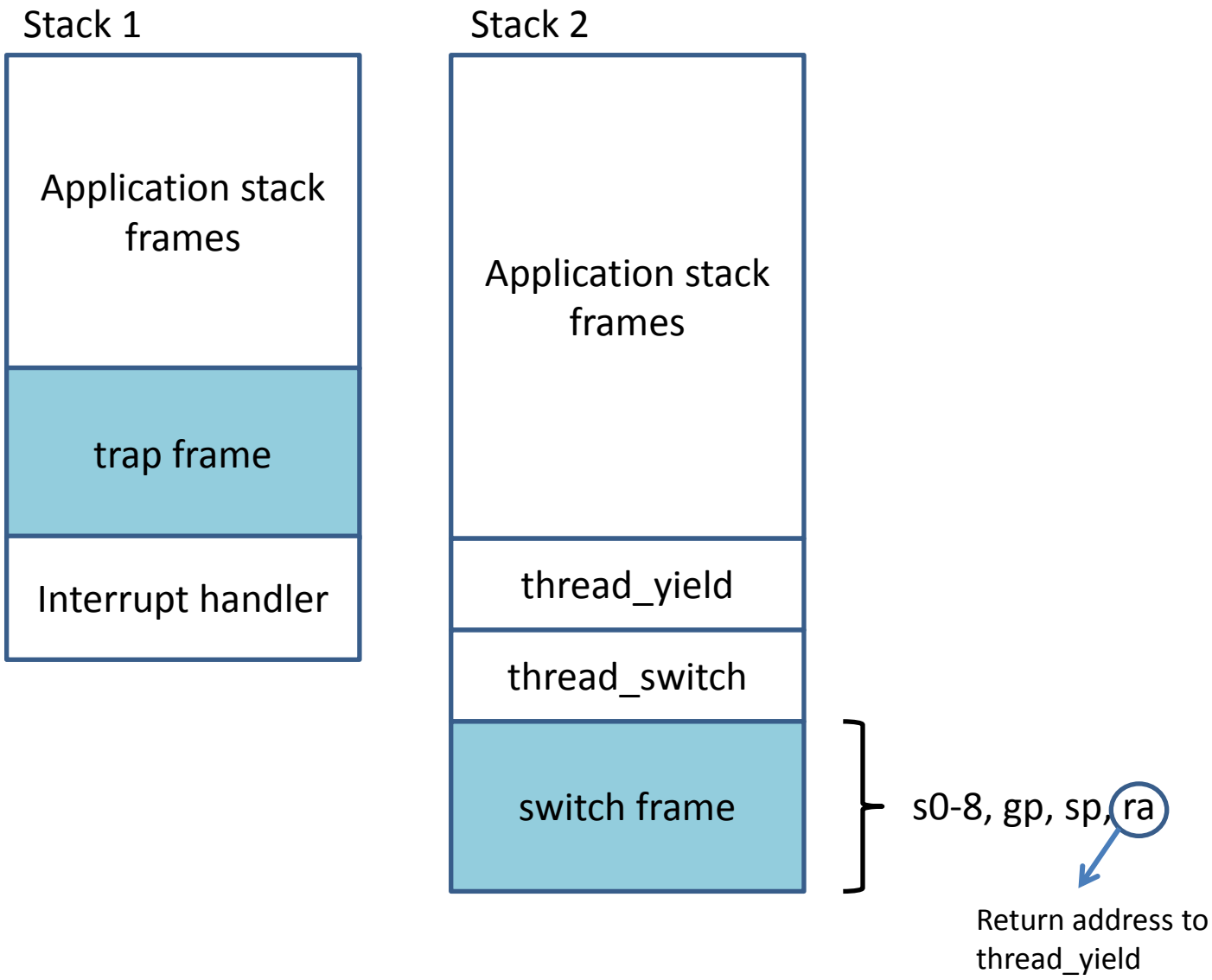
Transfers control to a fixed location in memory



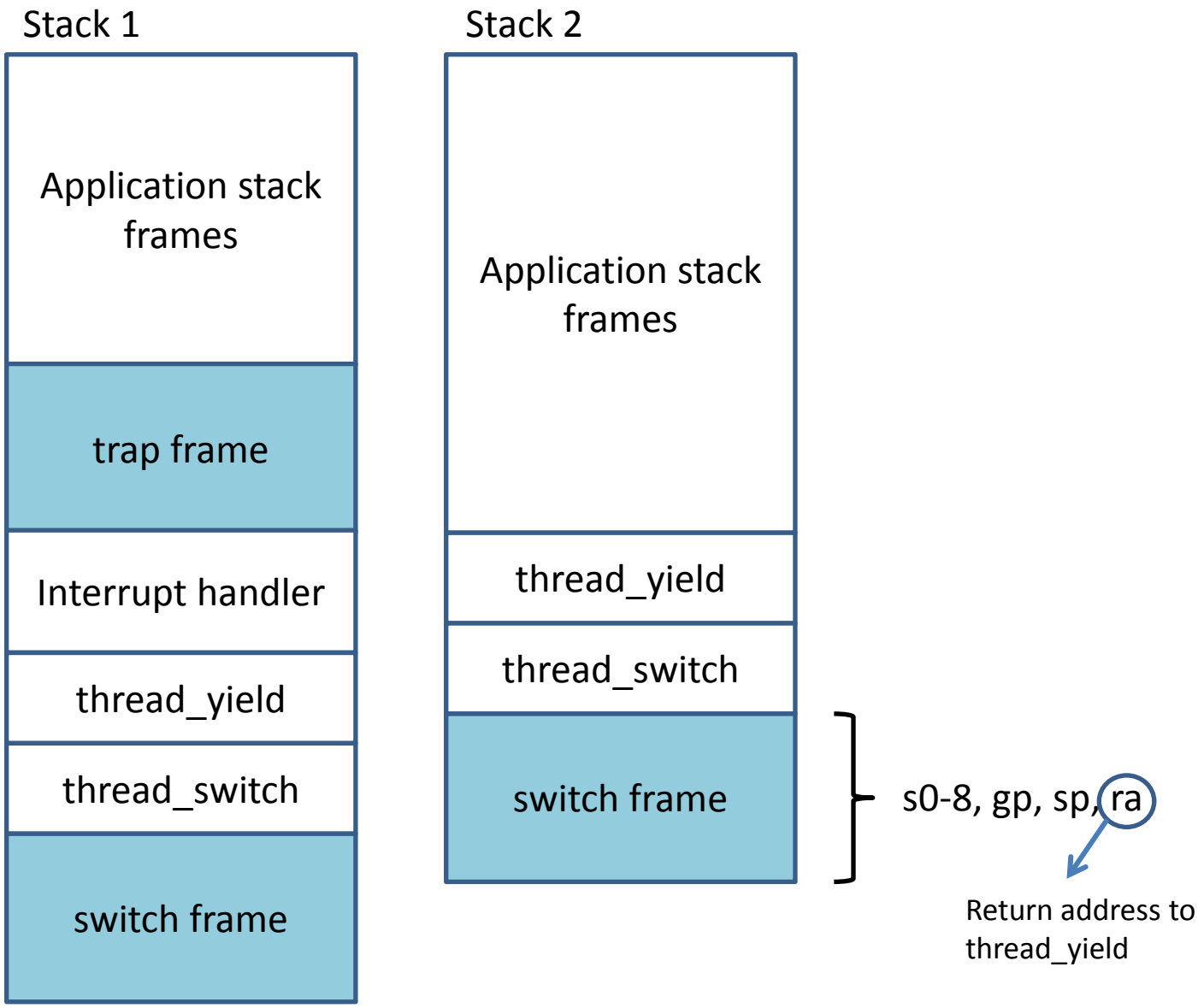
Save the current context into the trap frame. Must save every register.



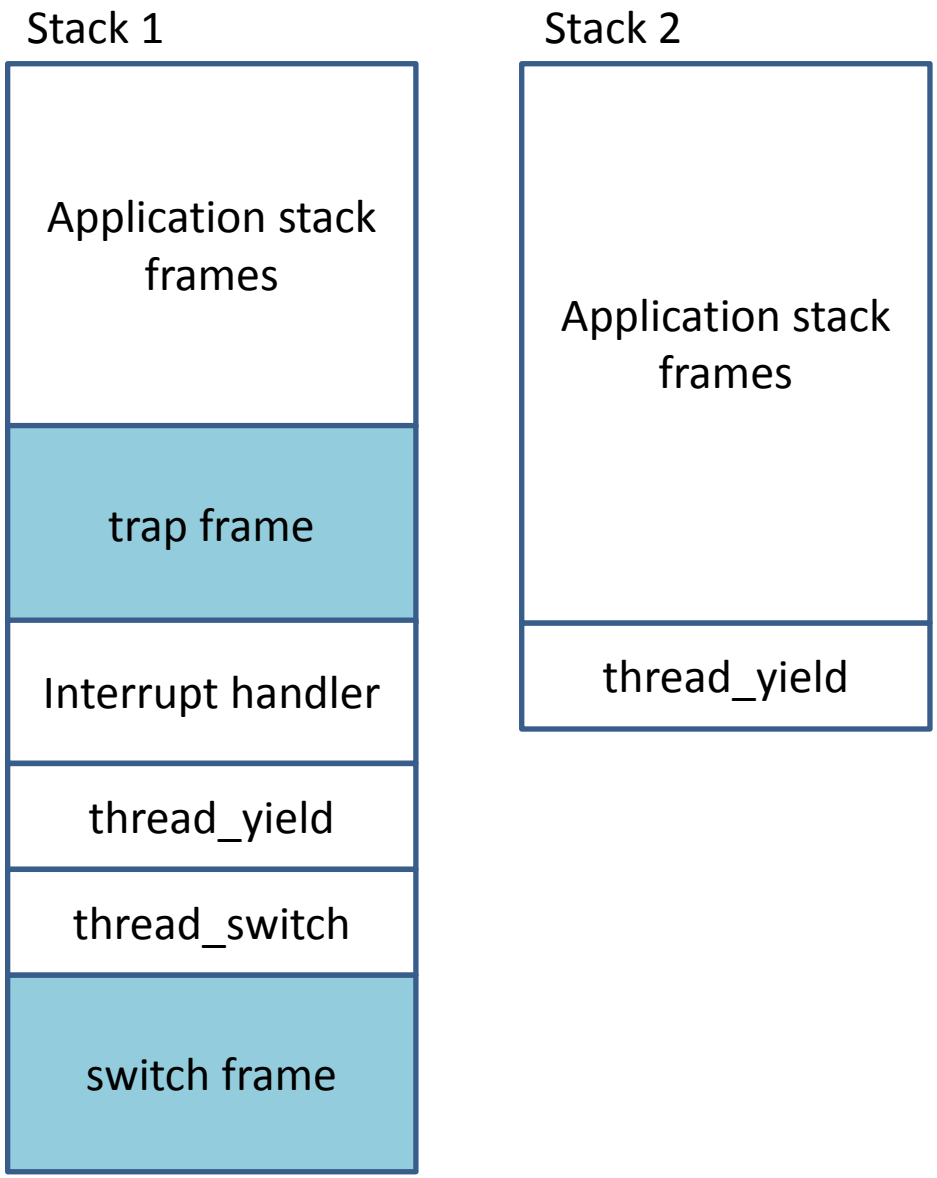
Call interrupt handler. Identify the device that raised the interrupt.



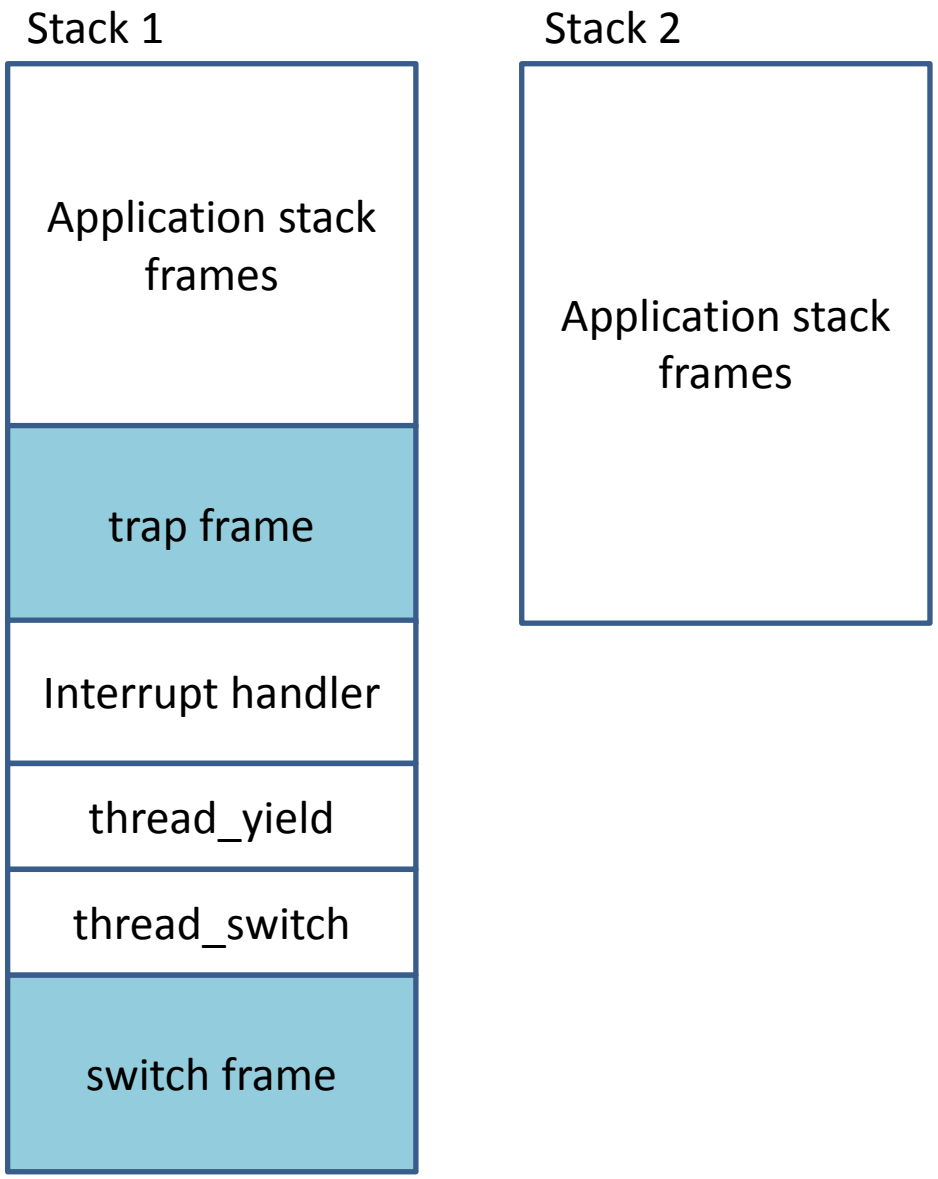
Running total counter equals to the quantum. Schedule thread 2 for execution. Perform context switch.



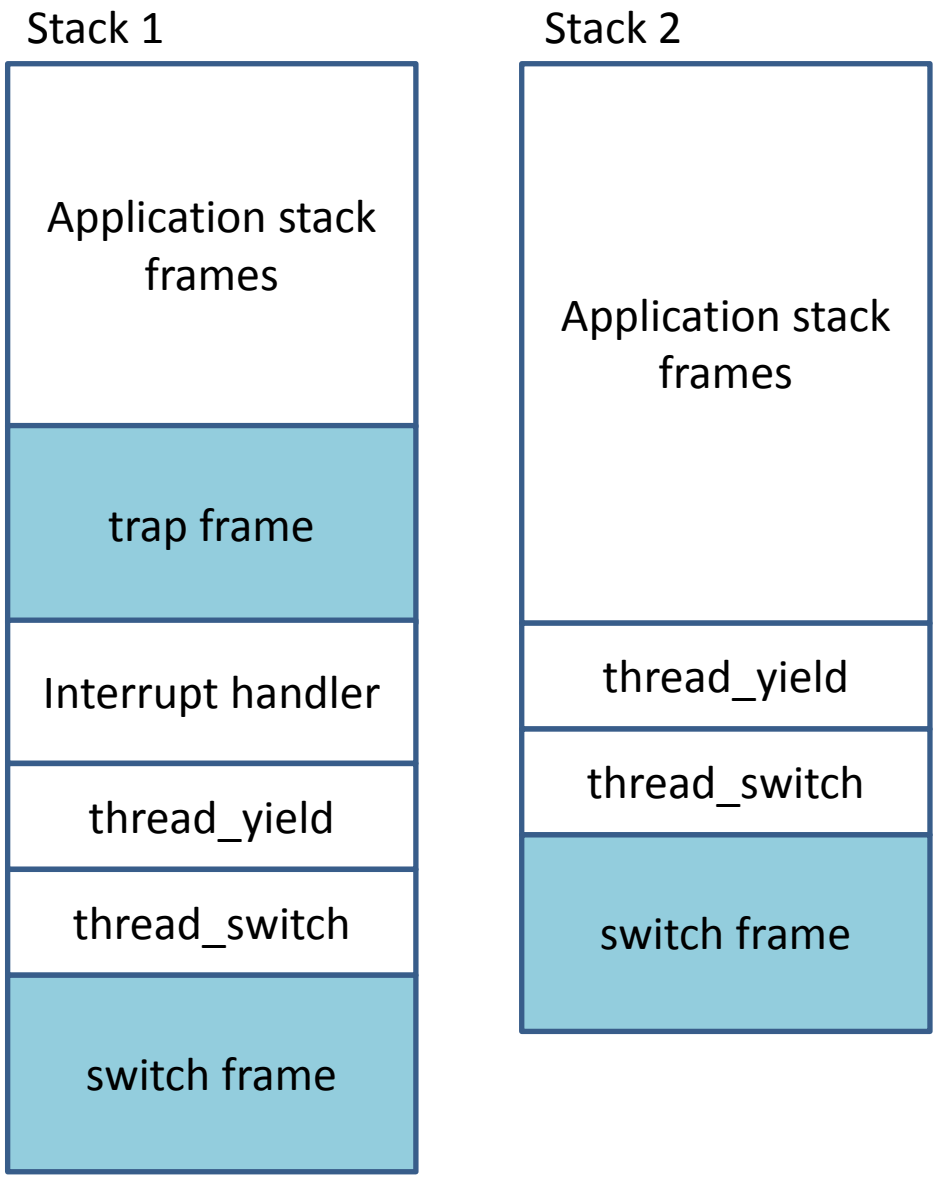
Restore the context from the switch frame. Returns to thread_yield using the value in return address (ra) register.



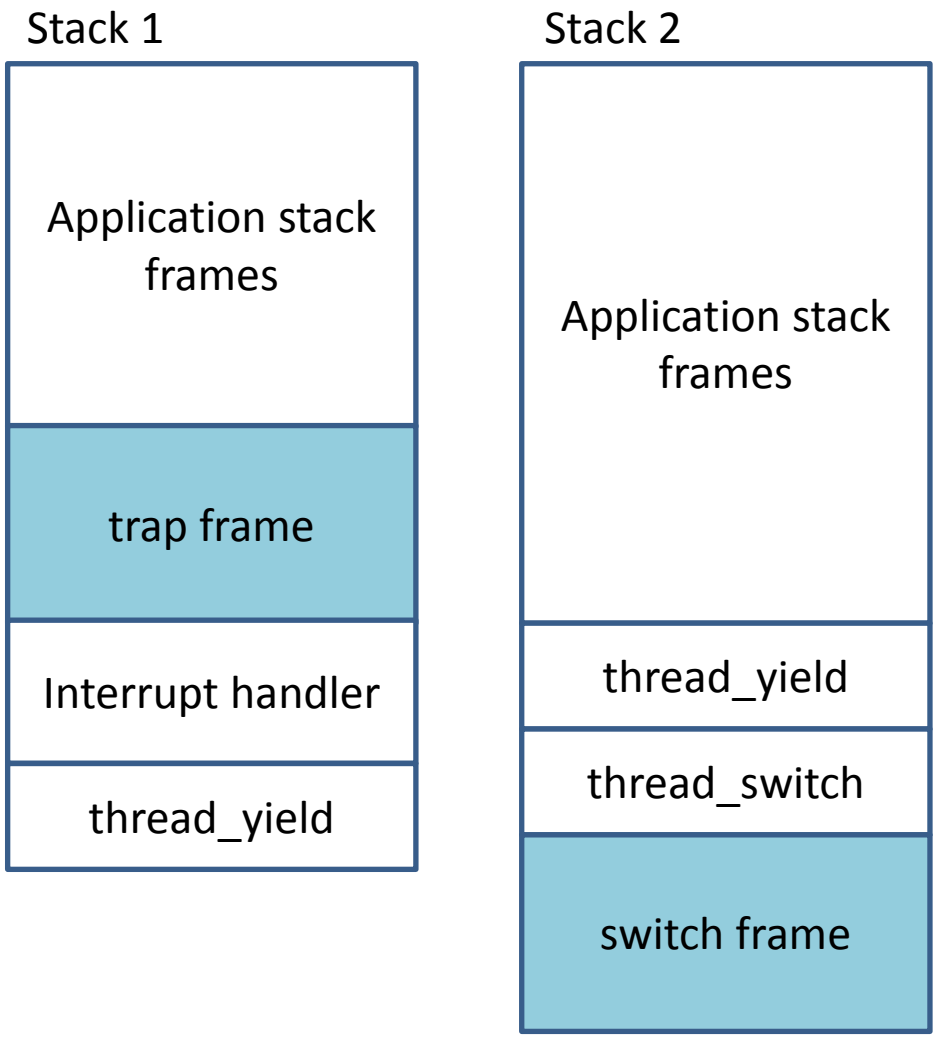
Return from thread_yield to continue executing the application.



Context switch back to thread 1.



Restore context from the switch frame and return on the ra register to continue in thread_yield.



After returning from the interrupt handler, restore the context in the trap frame and return to the application.