

Threads and Concurrency

key concepts

threads, concurrent execution, timesharing, context switch, interrupts, preemption

reading

Three Easy Pieces: Chapter 26 (Concurrency and Threads)

What is a Thread?

- Threads provide a way for programmers to express *concurrency* in a program.
- A normal *sequential program* consists of a single thread of execution.
- In threaded concurrent programs there are multiple threads of execution, all occurring at the same time.

OS/161 Threaded Concurrency Examples

- Key ideas from the examples:
 - A thread can create new threads using `thread_fork`
 - New threads start execution in a function specified as a parameter to `thread_fork`
 - The original thread (which called `thread_fork` and the new thread (which is created by the call to `thread_fork`) proceed concurrently, as two simultaneous sequential threads of execution.
 - All threads *share* access to the program's global variables and heap.
 - Each thread's function activations are *private* to that thread.

OS/161's Thread Interface

- create a new thread:

```
int thread_fork(  
    const char *name,           // name of new thread  
    struct proc *proc,         // thread's process  
    void (*func)                // new thread's function  
        (void *, unsigned long),  
    void *data1,                // function's first param  
    unsigned long data2         // function's second param  
);
```

- terminate the calling thread:

```
void thread_exit(void);
```

- volutarily yield execution:

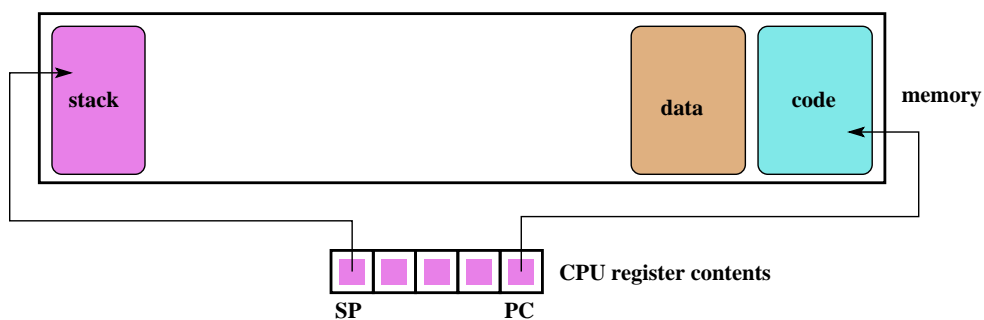
```
void thread_yield(void);
```

See `kern/include/thread.h`

Why Threads?

- **Reason #1:** parallelism exposed by threads enables parallel execution if the underlying hardware supports it.
 - programs can run faster
- **Reason #2:** parallelism exposed by threads enables better processor utilization
 - if one thread has to *block*, another may be able to run

Review: Sequential Program Execution



The Fetch/Execute Cycle

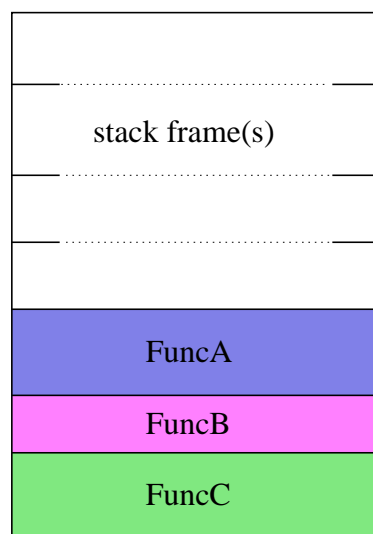
1. fetch instruction PC points to
2. decode and execute instruction
3. advance PC

MIPS Registers

num	name	use	num	name	use
0	z0	always zero	24-25	t8-t9	temps (caller-save)
1	at	assembler reserved	26-27	k0-k1	kernel temps
2	v0	return val/syscall #	28	gp	global pointer
3	v1	return value	29	sp	stack pointer
4-7	a0-a3	subroutine args	30	s8/fp	frame ptr (callee-save)
8-15	t0-t7	temps (caller-save)	31	ra	return addr (for jal)
16-23	s0-s7	saved (callee-save)			

See kern/arch/mips/include/kern/regdefs.h

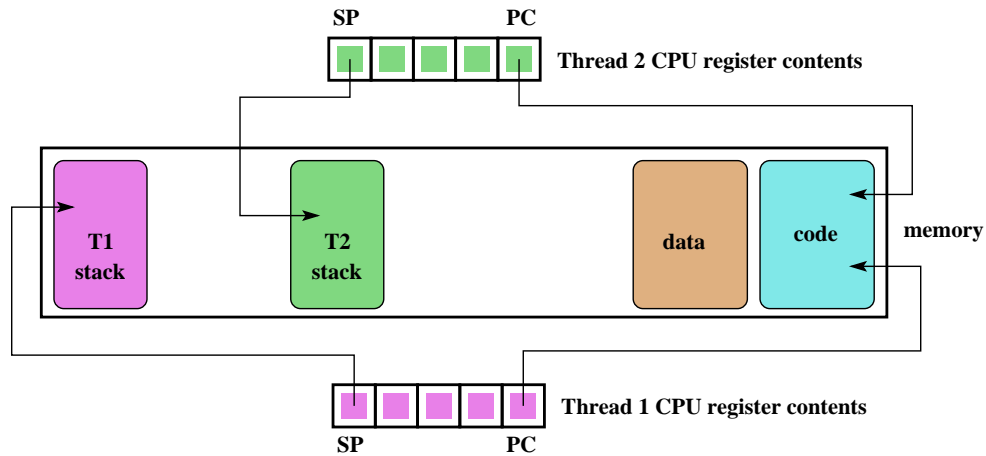
Review: The Stack



```
FuncA () {
    ...
    FuncB ();
    ...
}
```

```
FuncB () {
    ...
    FuncC ();
    ...
}
```

Concurrent Program Execution (Two Threads)



Conceptually, each thread executes sequentially using its private register contents and stack.

Implementing Concurrent Threads

- Option 1: multiple processors, multiple cores, hardware multithreading per core
 - P processors, C cores per processor, M multithreading degree per core \Rightarrow PCM threads can execute *simultaneously*
 - separate register set for each running thread, to hold its *execution context*
- Option 2: *timesharing*
 - multiple threads take turns on the same hardware
 - rapidly switch from thread to thread so that all make progress

In practice, both techniques can be combined.

Timesharing and Context Switches

- When timesharing, the switch from one thread to another is called a *context switch*
- What happens during a context switch:
 1. decide which thread will run next (scheduling)
 2. save register contents of current thread
 3. load register contents of next thread
- Thread context must be saved/restored carefully, since thread execution continuously changes the context

Context Switch on the MIPS (1 of 2)

```
/* See kern/arch/mips/thread/switch.S */

switchframe_switch:
    /* a0: address of switchframe pointer of old thread. */
    /* a1: address of switchframe pointer of new thread. */

    /* Allocate stack space for saving 10 registers. 10*4 = 40 */
    addi sp, sp, -40

    sw    ra, 36(sp) /* Save the registers */
    sw    gp, 32(sp)
    sw    s8, 28(sp)
    sw    s6, 24(sp)
    sw    s5, 20(sp)
    sw    s4, 16(sp)
    sw    s3, 12(sp)
    sw    s2, 8(sp)
    sw    s1, 4(sp)
    sw    s0, 0(sp)

    /* Store the old stack pointer in the old thread */
    sw    sp, 0(a0)
```

Context Switch on the MIPS (2 of 2)

```
/* Get the new stack pointer from the new thread */
lw    sp, 0(a1)
nop                    /* delay slot for load */

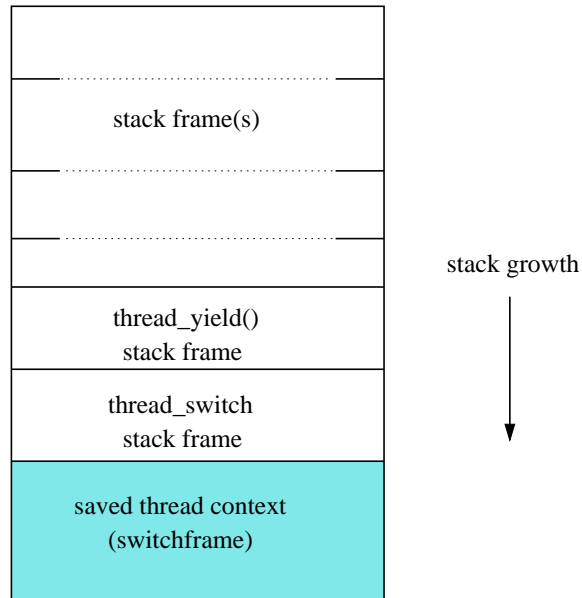
/* Now, restore the registers */
lw    s0, 0(sp)
lw    s1, 4(sp)
lw    s2, 8(sp)
lw    s3, 12(sp)
lw    s4, 16(sp)
lw    s5, 20(sp)
lw    s6, 24(sp)
lw    s8, 28(sp)
lw    gp, 32(sp)
lw    ra, 36(sp)
nop                    /* delay slot for load */

/* and return. */
j ra
addi sp, sp, 40        /* in delay slot */
.end switchframe_switch
```

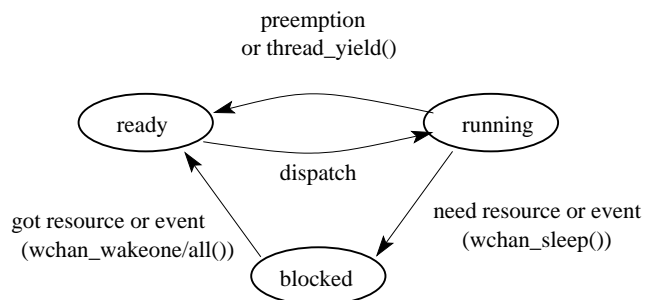
What Causes Context Switches?

- the running thread calls **thread_yield**
 - running thread *voluntarily* allows other threads to run
- the running thread calls **thread_exit**
 - running thread is terminated
- the running thread *blocks*, via a call to **wchan_sleep**
 - more on this later . . .
- the running thread is *preempted*
 - running thread *involuntarily* stops running

OS/161 Thread Stack after Voluntary Context Switch (`thread_yield()`)



Thread States



running: currently executing

ready: ready to execute

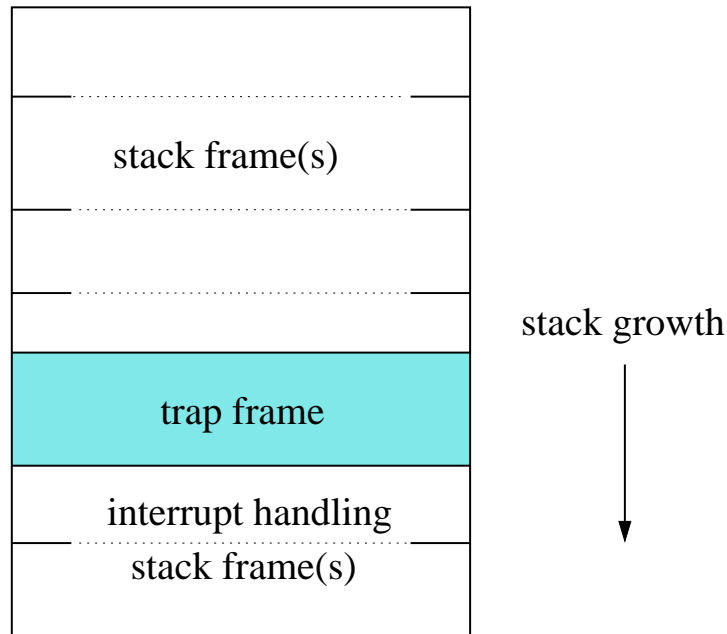
blocked: waiting for something, so not ready to execute.

Preemption

- without preemption, a running thread could potentially run forever, without yielding, blocking, or exiting
- *preemption* means forcing a running thread to stop running, so that another thread can have a chance
- to implement preemption, the thread library must have a means of “getting control” (causing thread library code to be executed) even though the running thread has not called a thread library function
- this is normally accomplished using *interrupts*

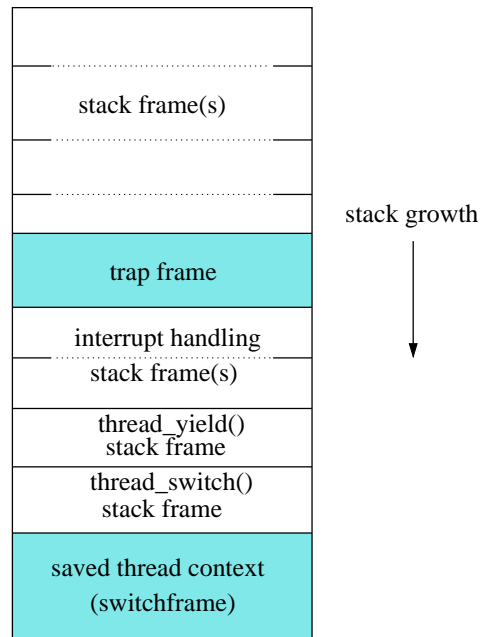
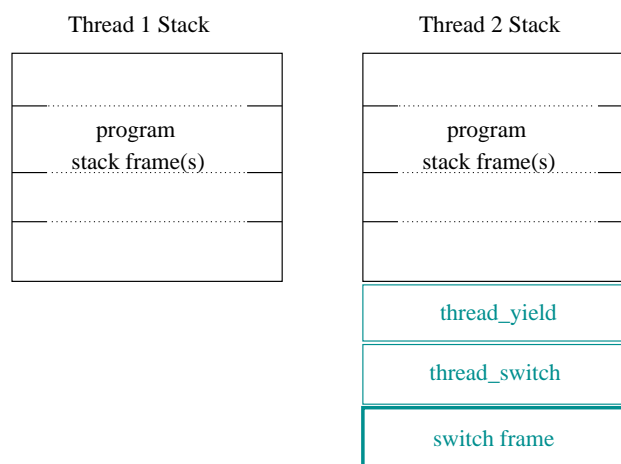
Review: Interrupts

- an interrupt is an event that occurs during the execution of a program
- interrupts are caused by system devices (hardware), e.g., a timer, a disk controller, a network interface
- when an interrupt occurs, the hardware automatically transfers control to a fixed location in memory
- at that memory location, the thread library places a procedure called an *interrupt handler*
- the interrupt handler normally:
 1. create a *trap frame* to record thread context at the time of the interrupt
 2. determines which device caused the interrupt and performs device-specific processing
 3. restores the saved thread context from the trap frame and resumes execution of the thread

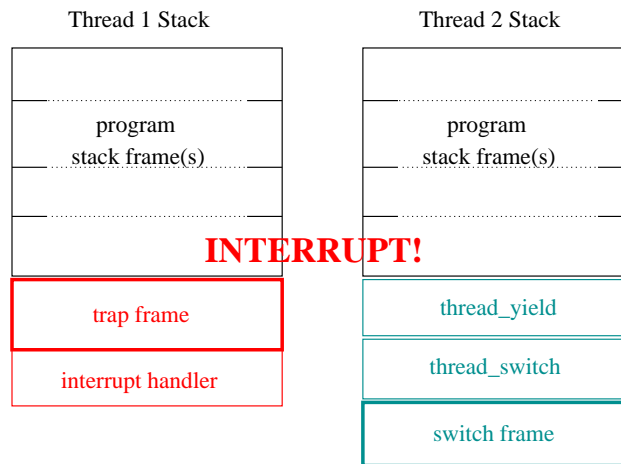
OS/161 Thread Stack after in Interrupt**Preemptive Scheduling**

- A preemptive scheduler imposes a limit, called the *scheduling quantum* on how long a thread can run before being preempted.
- The quantum is an *upper bound* on the amount of time that a thread can run. It may block or yield before its quantum has expired.
- Periodic timer interrupts allow running time to be tracked.
- If a thread has run too long, the timer interrupt handler preempts the thread by calling `thread_yield`.
- The preempted thread changes state from running to ready, and it is placed on the *ready queue*.

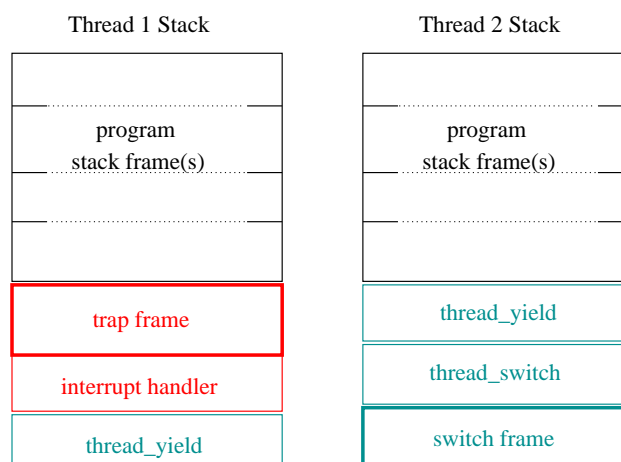
OS/161 threads use *preemptive round-robin scheduling*.

OS/161 Thread Stack after Preemption**Two-Thread Example (Part 1)**

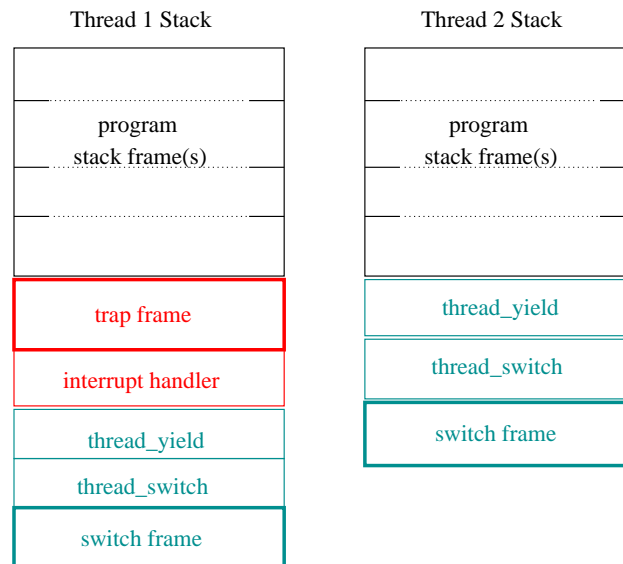
Thread 1 is running, thread two had previously yielded voluntarily.

Two-Thread Example (Part 2)

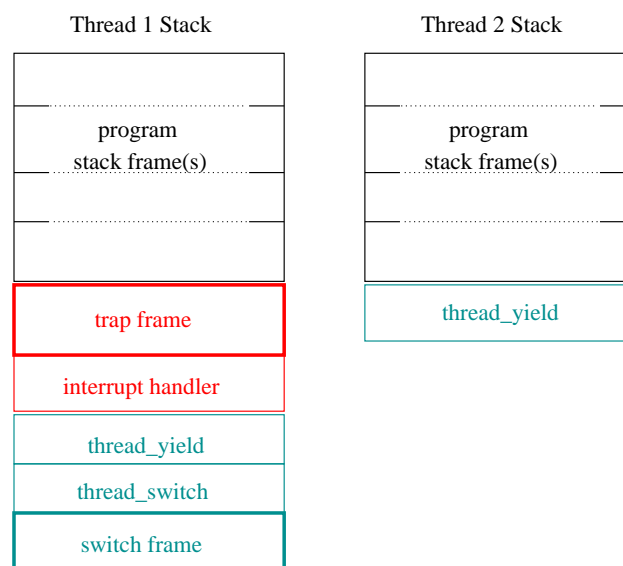
A time interrupt occurs! Interrupt handler runs.

Two-Thread Example (Part 3)

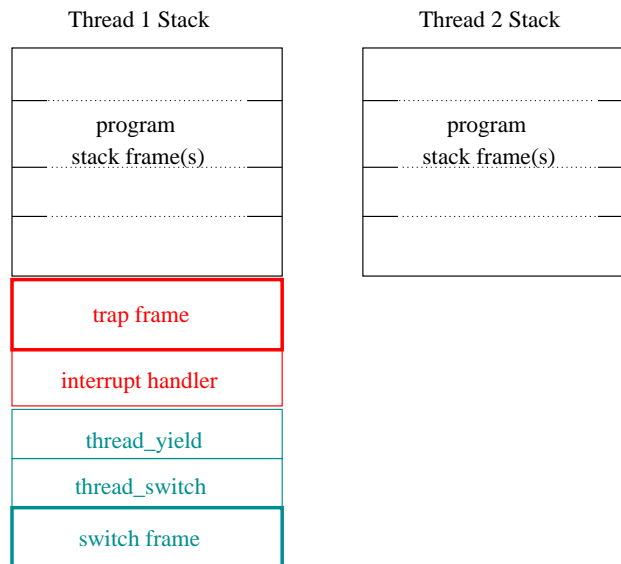
Interrupt handler decides Thread 1 quantum has expired.

Two-Thread Example (Part 4)

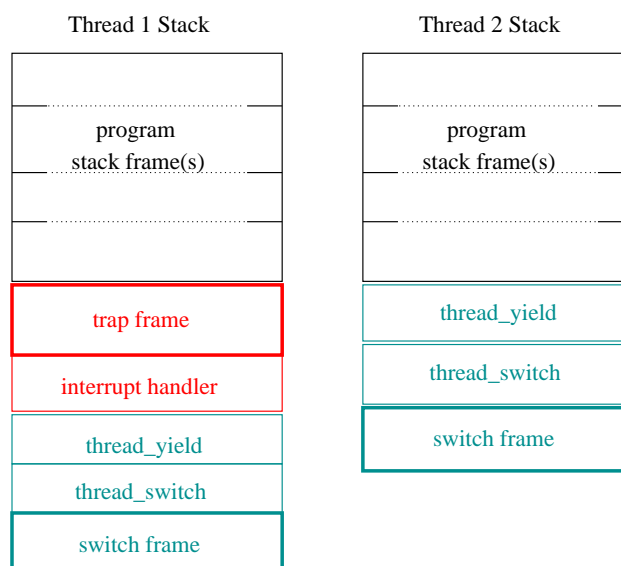
Scheduler chooses Thread 2 to run. Context switch.

Two-Thread Example (Part 5)

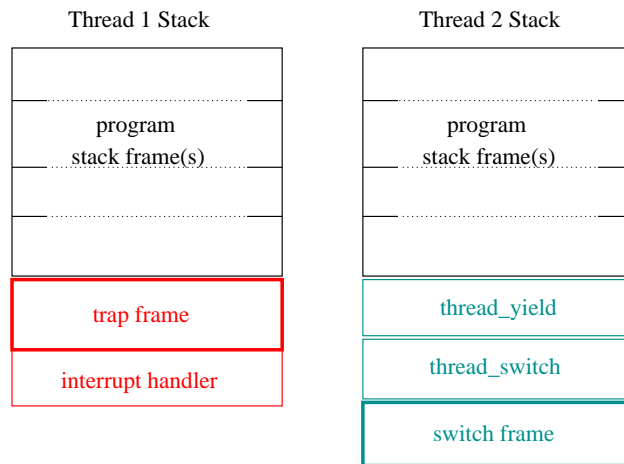
Thread 2 context is restored.

Two-Thread Example (Part 6)

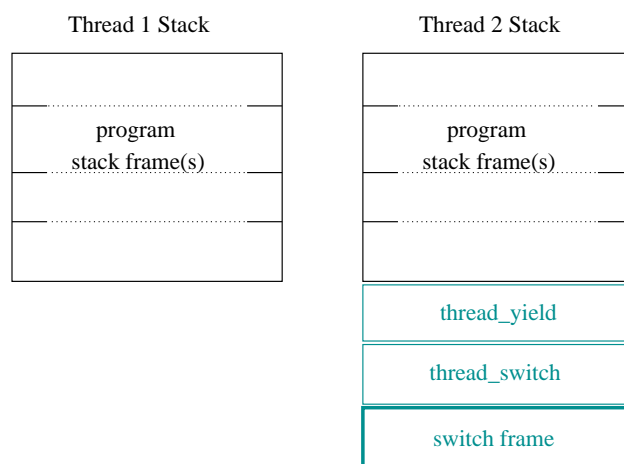
thread_yield finishes, Thread 2 program resumes.

Two-Thread Example (Part 7)

Later, Thread 2 yields again. Scheduler chooses Thread 1.

Two-Thread Example (Part 8)

Thread 1 context is restored, interrupt handler resumes.

Two-Thread Example (Part 9)

Interrupt handler restores state from trap frame and returns.
