

**University of Waterloo**  
**CS 462 — Formal Languages and Parsing**  
**Winter 2020**  
**Problem Set 7**

*Distributed Friday, February 28 2020. Due Friday March 6 2020 at 5 PM. Submit via LEARN.*

1. [10 marks] Consider the positive squares, represented in base 6. Here are the first few elements of this language:

$$L = \{1, 4, 13, 24, 41, 100, 121, 144, 213, 244, \dots\}.$$

What are the minimal elements of  $L$  under the subsequence order? Prove your answer.

2. [10 marks] Call a morphism  $h : \Sigma^* \rightarrow \Sigma^*$  non-erasing if  $h(a) \neq \epsilon$  for all  $a \in \Sigma$ .

Let the domain be  $S = \{\epsilon\} \cup 0\{0, 1\}^*$ . These are the binary strings that do not begin with a 1.

Consider the following order defined on  $S$ : for strings  $x, y \in S$  we say  $x \leq y$  if there is some non-erasing morphism  $h : \Sigma^* \rightarrow \Sigma^*$  such  $h(x) = y$ .

Examples:  $010 \leq 000$  because we can take  $h(0) = h(1) = 0$ .

$010 \leq 1001$  because we can take  $h(0) = 1, h(1) = 00$ .

But  $0110$  is not  $\leq 0101$ , because whatever we decide to map 1 to must be doubled in its image under  $h$ . So it must be that  $h(1) = 01$ , the only doubled string in  $0101$ . But since  $h(0)$  is nonempty, this says  $h(0110)$  must be of length at least 6, contradiction.

(a) Show that  $\leq$  is a partial order on  $S$ .

(b) Recall that an *antichain* is a set of pairwise incomparable elements. Show that  $\leq$  has infinite antichains.

3. [10 marks] If  $L \subseteq \Sigma^*$  is a CFL, need

$$\{x \in \Sigma^* : x^* \subseteq L\}$$

be a CFL? Prove or give a counterexample.