

## ASSIGNMENT 4

DO NOT COPY. ACKNOWLEDGE YOUR SOURCES.

1. [10 marks] Prove that  $\text{RP} \cap \text{co-RP} \subseteq \text{ZPP}$ . See Lecture 8 notes for definitions of these classes. Recall that I gave the idea of this proof in class: assume you have an RP algorithm  $A$  and a co-RP algorithm  $A'$ . By running these (repeatedly) build a ZPP algorithm. Include probability details in your solution.
2. [10 marks] Two multisets of numbers are equal if they contain the same numbers with the same multiplicities, e.g.,  $\{1, 2, 1\}$  and  $\{2, 1, 1\}$  are equal, but  $\{1, 2, 2\}$  is not. We can test if two multisets of size  $n$  are equal by sorting in  $O(n \log n)$  or by hashing, but here is another method.

For a multiset  $S = \{e_1, e_2, \dots, e_n\}$  define  $P(S) = (x - e_1)(x - e_2) \cdots (x - e_n)$ , a polynomial in one variable  $x$ .

- (a) [3 marks] For multisets  $S_1$  and  $S_2$  prove that  $S_1 = S_2$  iff  $P(S_1) \equiv P(S_2)$ . (You may quote results from linear algebra without proof.)
- (b) [7 marks] Based on this, and on the Schwartz-Zippel result from class, suggest a randomized Monte Carlo algorithm to test if  $S_1 = S_2$ . Your algorithm should have error probability at most  $1/2$ . Be clear about the range of your random numbers. On which side are the one-sided errors? Analyze the run-time of your algorithm, assuming arithmetic operations take constant time.

For practicality, your computations should be done modulo a small prime, but you do not need to analyze how large a prime is needed.